

Suppose A, A' are presentation matrices

for the R -module M . Then

$$\begin{array}{ccccc}
 & \xrightarrow{F} & A & \xrightarrow{\pi} & M \\
 F \xrightarrow{\quad} & \xrightarrow{E} & \text{---} & \xrightarrow{\quad} & AD = B'B - \text{Id} \\
 \downarrow C & & \downarrow C' & & \\
 C \downarrow F & \xrightarrow{A'} & E' & \xrightarrow{\quad} & BA = A'C \\
 & \downarrow B & \downarrow B' & \parallel & B'A' = AC' \\
 & \xrightarrow{F'} & \xrightarrow{E} & \xrightarrow{\quad} & M \xrightarrow{\quad} M
 \end{array}$$

Using that E, E', F, F' are free, we construct B, B', C, C' s.t. the diagram commutes.

$$\pi \circ (B'B - \text{Id}) = 0 \Rightarrow \exists D$$

We'd like to prove $A \sim A'$

First, up to adding and killing generators,

we may suppose A, A' with the same number of rows.

(it was described by $A \sim \left(\begin{array}{c|c} A & 0 \\ \hline 0 & \text{Id} \end{array} \right)$)

$$A \sim \left(\begin{array}{c|c} A & 0 \\ \hline 0 & \text{Id} \end{array} \right) \sim \left(\begin{array}{c|c} A & B' \\ \hline 0 & \text{Id} \end{array} \right) \sim \left(\begin{array}{c|c|c} A & B' & 0 \\ \hline 0 & \text{Id} & 0 \end{array} \right)$$

$$\sim \left(\begin{array}{c|c|c} A & B' & B'A' \\ \hline 0 & \text{Id} & A' \end{array} \right) = \left(\begin{array}{c|c|c} A & B' & AC' \\ \hline 0 & \text{Id} & A' \end{array} \right) \sim$$

$$\sim \left(\begin{array}{c|c|c} A & B' & 0 \\ \hline 0 & \text{Id} & A' \end{array} \right) \sim \left(\begin{array}{c|c|c|c} A & B' & 0 & B'B \\ \hline 0 & \text{Id} & A' & B \end{array} \right)$$

$$= \left(\begin{array}{c|c|c|c|c} A & B' & 0 & AD + Id \\ \hline 0 & \text{Id} & A' & B \end{array} \right) \sim \left(\begin{array}{c|c|c|c} A & B' & 0 & Id \\ \hline 0 & \text{Id} & A' & B \end{array} \right)$$

$$A' \sim \left(\begin{array}{c|c|c|c} A' & B & 0 & Id \\ \hline 0 & \text{Id} & A & B' \end{array} \right)$$

obtain one
from the other
via permutation
of rows and columns

This proves that $A \sim A'$.

Definition : M finitely presented R -module

with presentation matrix A with m rows

(which means m generators for M). The

n -th elementary ideal $E_n(M)$

is the ideal in R generated by

$(m-n+1) \times (m-n+1)$ minors of A .

It is well-defined thanks to the

Theorem stating that all presentation matrices are equivalent (the only non-trivial fact being that A and $\left(\begin{smallmatrix} A & 0 \\ 0 & 1 \end{smallmatrix} \right)$ generate the same

ideals; we replace). By replace,

a $(d+1) \times (d+1)$ minor is generated by

$d \times d$ minors. $\Rightarrow E_n(M) \subseteq E_{n+1}(M) \quad \forall n.$

We agree that $E_n(M) = R$ if $n > m$

$E_n(M) = 0$ if $n \leq 0.$

Exercise: $\mathbb{Z}[t, t^{-1}]$ is not a P.I.D., but

it is a U.F.D. so $E_n(M)$ is not

principal in general, but it has a

G.C.D. (greatest common divisor), which

is named the n -th Alexander polynomial

of M (or L , when $M = H_1(\tilde{E}_\infty)$)

$E = S^3, L).$

② The invertibles of $\mathbb{Z}[t, t^{-1}]$ are

$\pm t^n$, $n \in \mathbb{Z}$. Hence Alexander

polynomials are well-defined only up

to multiplication by $\pm t^n$

Definition: "The" Alexander polynomial is the first Alexander polynomial.

Notation: $f, g \in \mathbb{Z}[t, t^{-1}]$. We write

$$f \doteq g \iff f = \pm t^n g.$$

Example: K the unknot. $S^3 \setminus K \cong$

$$\cong \mathbb{R}^3 \setminus \{\text{straight line}\} \cong S^1 \times \mathbb{R}^+ \times \mathbb{R},$$

$$\text{hence } \widetilde{E}_\infty = \widetilde{E} \text{ (universal covering)} = \mathbb{R} \times \mathbb{R}^+ \times \mathbb{R}$$

$$\cong \mathbb{R}^3 \Rightarrow H_1(\widetilde{E}_\infty) = 0, \text{ which is presented}$$

e.g. by the matrix (1). Hence $E_1(K) = \mathbb{R}$

and the first Alexander polynomial is

$$\Delta_K(t) = 1.$$

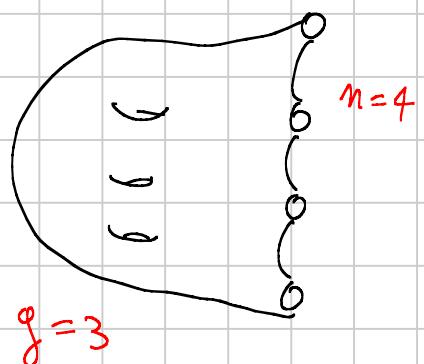
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notation for the 1-st Alexander polynomial

Problem: How to compute Δ_L in general?

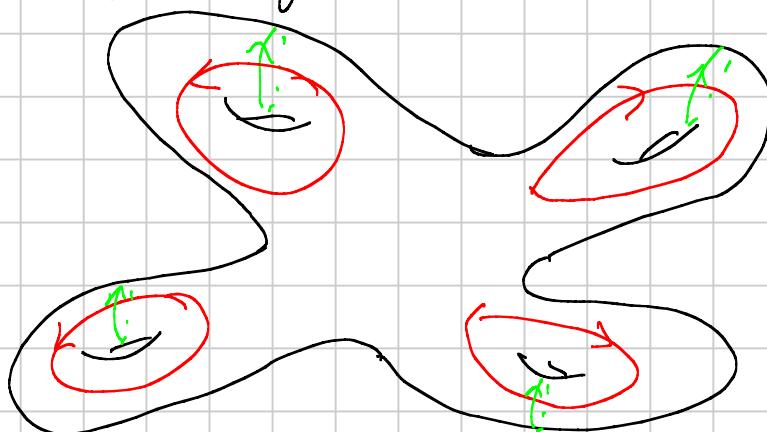
1-st Homology of surfaces

Orientable surface is $\sum_{g,m} =$

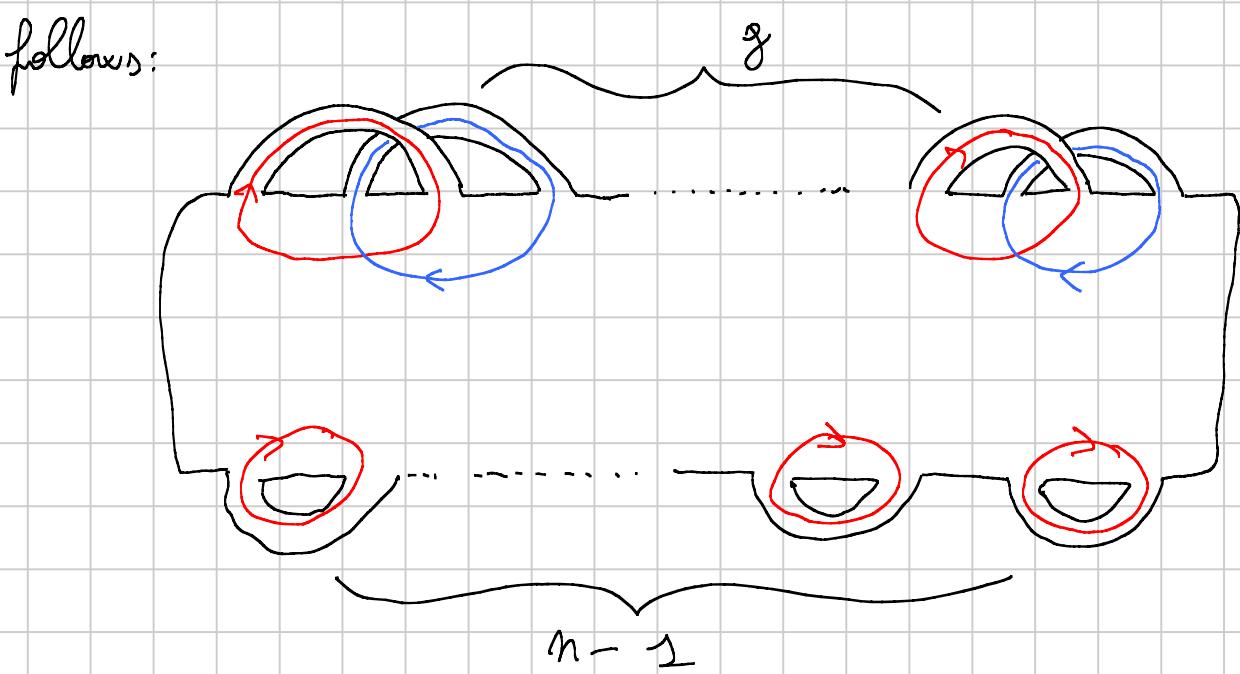


If $m=0$, $H_1(\Sigma_{g,0}) = H_1(\Sigma_g) \cong \mathbb{Z}^{2g}$

with the following basis



If $m > 0$ (i.e. $\partial \Sigma_{g,m} \neq \emptyset$) the $\Sigma_{g,m}$ is as follows:



$$H_1(\Sigma_{g,n}) \cong \mathbb{Z}^{2g+n-1} \text{ with the basis above}$$

We come back to our setting: $L \subseteq S^3$

and we look for a presentation of $H_1(\tilde{E}_\infty)$,

$$E = S^3 \setminus L.$$

We also fix a Seifert surface S for L .

$$n = \# \text{ components of } L, \quad g = \text{genus}(S).$$

$$\text{Prop.: } H_1(S) = \mathbb{Z}^{2g+n-1} \cong H_1(S^3 \setminus S)$$

and \exists a bilinear $\beta: H_1(S^3 \setminus S) \times H_1(S) \rightarrow \mathbb{Z}$

s.t.

$$\beta([c], [d]) = \text{lk}(c, d)$$

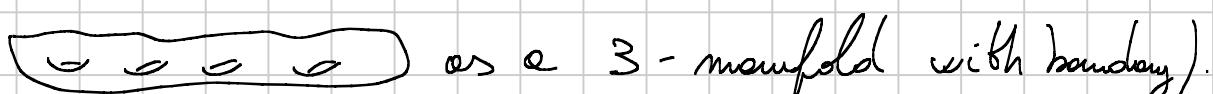
\forall simple closed oriented loops $c \subseteq S^3 \setminus S, d \subseteq S$

$$\text{Proof: } V = \overline{N(S)}, \quad V' = \overline{S^3 \setminus V}$$

By representing S as a disk with (knotted)

bonds, we understand $V = \text{handlebody of}$

genus $2g+n-1$ ($\text{handlebody of genus } g =$



Of course, V is a **knotted handlebody**,

with one (knotted) solid cylinder coming from each (knotted) band of S . On $\partial V = \sum_{i=2g+m-1}^n$

we take the standard basis for H_1

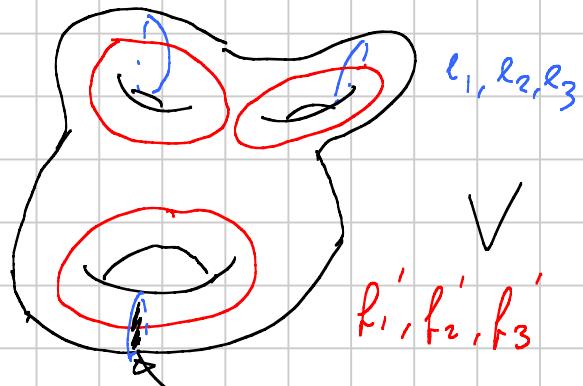
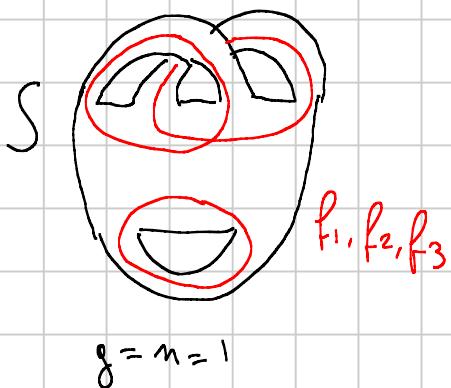
which can be constructed as follows:

f_1, \dots, f_{2g+m-1} "standard" basis for $H_1(S)$

$f'_1, \dots, f'_{2g+m-1} \in H_1(\partial V)$ s.t. $f'_i = f_i$ in $H_1(V)$

$e_1, \dots, e_{2g+m-1} \in H_1(\partial V)$ s.t. e_i bounds a disk in V

"dual" to f_i :



From Mayer-Vietoris

$$H_2(S^3) \rightarrow H_1(\partial V) \xrightarrow{\psi} H_1(V) \oplus H_1(V') \rightarrow H_1(S^3)$$

$$\begin{matrix} \text{||} & & \text{||} \\ 0 & \longrightarrow & (f_i, ?) & 0 \end{matrix}$$

$$e_i \longrightarrow (0, ?)$$

ψ is an isomorphism \Rightarrow the e_i define a

basis of $H_1(V') \cong H_1(S^3 \setminus S)$.

Now define an orientation on f_i, e_i

s.t. $\text{lk}(f_i, e_i) = +1$ (recall that

e_i bounds a disc intersecting f_i transversely once).

(This is not bothered by knottedness of V, V').

Since this disc is disjoint from f_j , $j \neq i$,

$\text{lk}(f_j, e_i) = \delta_{ij}$. Now define

$$\beta(\sum c_i [e_i], \sum b_j [f_j]) = \sum c_i b_i$$

I now claim that $\beta([c], [d]) = \text{lk}(c, d)$.

$$\forall i, \text{lk}(c, f_i) = [c] \text{ in } H_1(S^3 \setminus f_i) =$$

$$= [\sum c_k [e_k]] = \sum c_k \text{lk}(e_k, f_i) = c_i$$

(I am assuming $[c] = \sum c_k [e_k]$, $[d] = \sum b_k [f_k]$).

$$\text{Thus } \text{lk}(c, d) = [d] \text{ in } H_1(S^3 \setminus c) =$$

$$= [\sum b_k [f_k]] = \sum b_k \text{lk}(f_k, c) = \sum b_k c_k =$$

$$= \beta([c], [d]).$$

Facts: From elementary linear algebra

(over \mathbb{Z}), for every basis f_1, \dots, f_{2g+m-1}

of $H_1(S)$ there exists a dual basis e_1, \dots, e_{2g+m-1}

s.t. $\beta(f_i, e_j) = \delta_{ij}$.

Remember S is oriented $\Rightarrow N(S) \cong S \times [-1, 1]$

hence we have embeddings $i^+: S \rightarrow S^3$ $i^-: S \rightarrow S^3$

$i^\pm(x) = (x, \pm 1)$. We define

$$\alpha: H_1(S) \times H_1(S) \rightarrow \mathbb{Z}$$

$$\alpha(x, y) = \beta(x^-, y) - \beta(y^+, x)$$

where $x^\pm = i_\ast^\pm(x)$, $y^\pm = i_\ast^\pm(y)$

clear from a geometric
point of view

$$\beta: H_1(S^3, S) \times H_1(S) \rightarrow \mathbb{Z}$$

Remark: Fix dual basis f_i, e_i of $H_1(S)$, $H_1(S^3, S)$. Let A be the matrix representing α w.r.t. the f_i , i.e.

$$A_{ij} = \alpha(f_i, f_j) = \beta(f_i^-, f_j^-).$$

Then $f_i^- = \sum_j A_{ij} e_j$, $f_i^+ = \sum_j A_{ji} e_j$

In fact, f_i^- and $\sum_j A_{ij} e_j$ have the same pairings with each f_k via β , hence they are equal.

The same for f_i^+ .

* Theorem : A representing α w.r.t. a basis of $H_1(S)$. Then $tA - {}^t\bar{A}$ is a presentation matrix for $H_1(\tilde{E}_\infty)$.